

Experiences
in the
Research and Development
of a
Non-clausal SAT Solver

John Franco, Sean Weaver

ECECS, University of Cincinnati

United States Department of Defense

Abridged Sample Input (from CMU webpage)

```
MODULE module-name-changed
INPUT
  ID_EX_RegWrite, ID_EX_MemToReg, _Taken_Branch_1_1, EX_MEM_Jump, ...
OUTPUT _temp_1252;
STRUCTURE
  _squash_1_1 = or(_Taken_Branch_1_1, EX_MEM_Jump);
  _squash_bar_1_1 = not(_squash_1_1);
  _EX_Jump_1_1 = and(_squash_bar_1_1, ID_EX_Jump);
  _Taken_Branch_9_1 = and(_squash_bar_1_1, ID_EX_Branch, TakeBranchALU_0);
  _Reg2Used_1_1 = or(IF_ID_UseData2, IF_ID_Branch, IF_ID_MemWrite, IF_ID_MemToReg);
  _temp_967 = and(_Reg2Used_1_1, e_2_1);
  ...
  _temp_976 = ite(_temp_969, IF_ID_Jump, Jump_0);
  ...
  _temp_1249 = and(_temp_1038, _temp_1066, _temp_1072, _temp_1189, _temp_1246);
  ...
  true_value = new_int_leaf(1);
  are_equal(_temp_1252, true_value); % 1
ENDMODULE
```

Courtesy M.N. Velev, Superscalar Suite 1.0. Available from: <http://www.ece.cmu.edu/~mvelev>.

Translation to CNF

Expression: $v_3 = ite(v_0, v_1, v_2)$; What Heuristics Like:

Karnaugh Map:

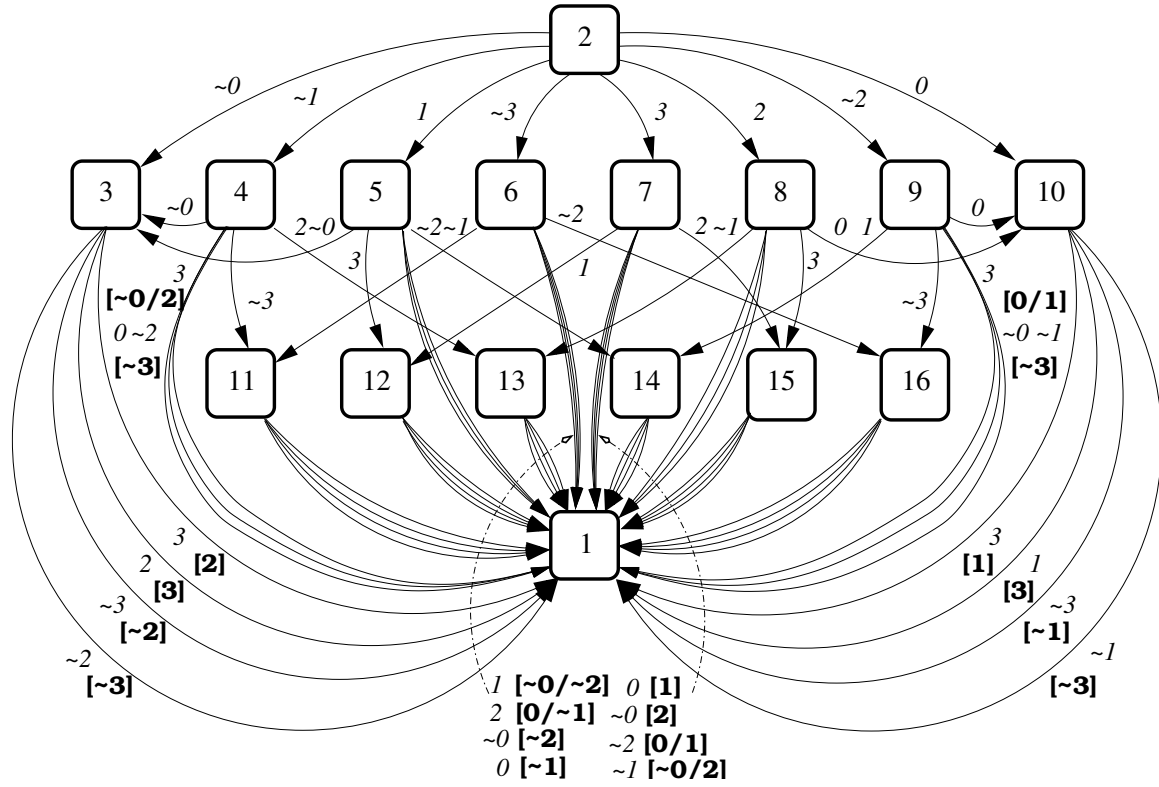
	00	01	11	10
00	1	0	1	0
01	1	0	1	0
11	0	1	1	0
10	1	0	0	1

CNF:

$$\begin{aligned}
 &(v_0 \vee v_2 \vee \bar{v}_3) \\
 &(v_0 \vee \bar{v}_2 \vee v_3) \\
 &(\bar{v}_0 \vee \bar{v}_1 \vee v_3) \\
 &(\bar{v}_0 \vee v_1 \vee \bar{v}_3)
 \end{aligned}$$

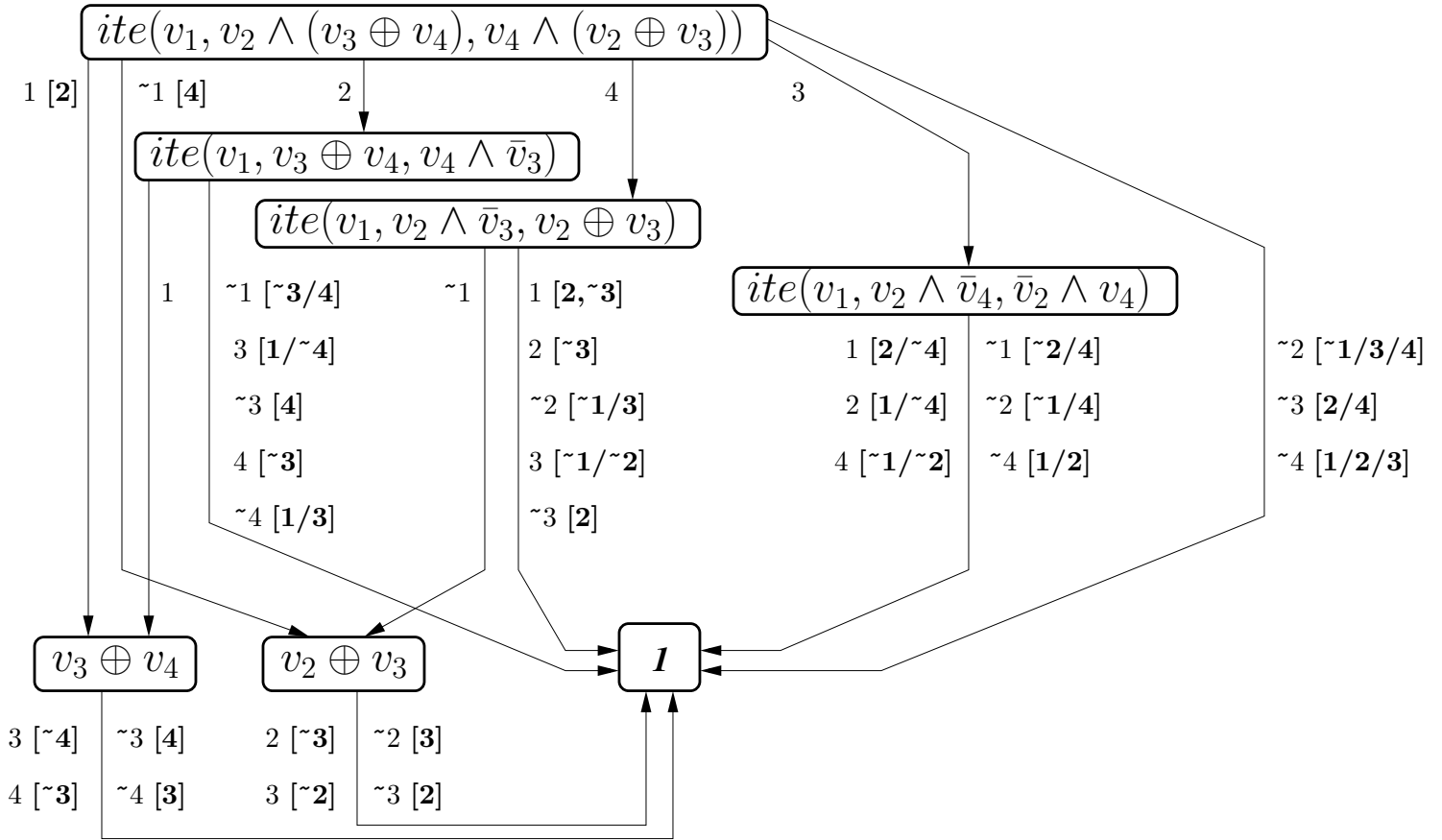
$$\begin{aligned}
 &(v_0 \wedge v_1) \rightarrow v_3 & (v_0 \wedge v_3) \rightarrow v_1 \\
 &(\bar{v}_0 \wedge v_2) \rightarrow v_3 & (\bar{v}_0 \wedge v_3) \rightarrow v_2 \\
 &(v_0 \wedge \bar{v}_1) \rightarrow \bar{v}_3 & (v_0 \wedge \bar{v}_3) \rightarrow \bar{v}_1 \\
 &(\bar{v}_0 \wedge \bar{v}_2) \rightarrow \bar{v}_3 & (\bar{v}_0 \wedge \bar{v}_3) \rightarrow \bar{v}_2 \\
 &(v_1 \wedge v_2) \rightarrow v_3 & (\bar{v}_1 \wedge \bar{v}_2) \rightarrow \bar{v}_3 \\
 &(v_1 \wedge \bar{v}_3) \rightarrow \bar{v}_0, \bar{v}_2 & (v_2 \wedge \bar{v}_3) \rightarrow v_0, \bar{v}_1 \\
 &(\bar{v}_1 \wedge v_3) \rightarrow \bar{v}_0, v_2 & (\bar{v}_2 \wedge v_3) \rightarrow v_0, v_1 \\
 &(v_1 \wedge \bar{v}_2) \rightarrow v_0 = v_3 \\
 &(\bar{v}_1 \wedge v_2) \rightarrow \bar{v}_0 = v_3
 \end{aligned}$$

State Machine Used to Represent Functions



States of a SMURF representing $v_3 = ite(v_0, v_1, v_2)$

SMURFs Don't Have to be Big



Locally Skewed, Globally Balanced Heuristics

Weight of terminal state is 0

If state s has p successors $\{s_1, s_2, \dots, s_p\}$, weight of s is

$$\frac{\sum_{i=1}^p \text{weightOf}(s_i) + \text{numberInferencesMadeEnrouteTo}(s_i)}{K * p}$$

Every state transition gets a weight:

inferences to destination state plus weight of that state

Every literal gets a score:

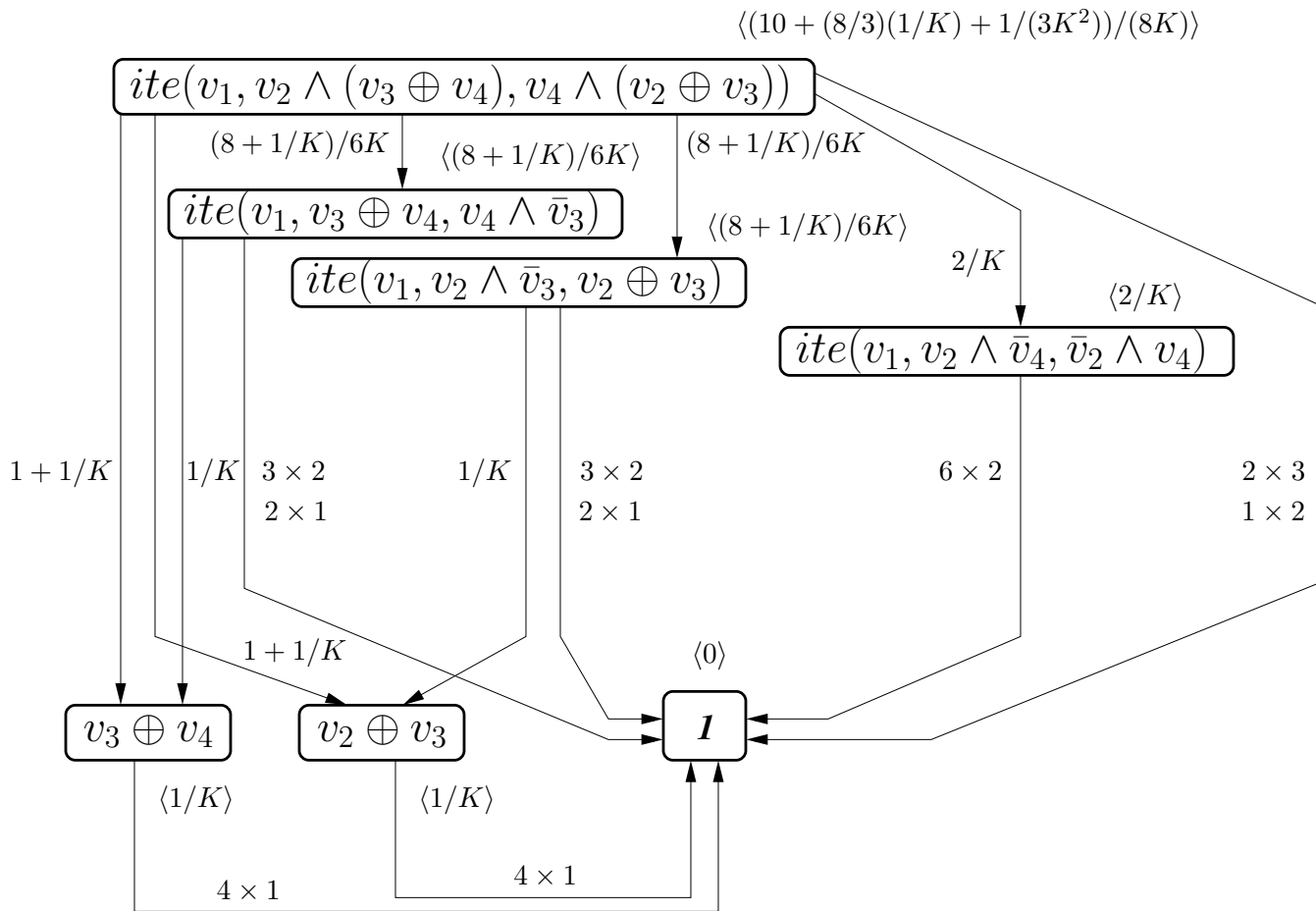
Sum of transition weights for that literal across SMURFs

Every variable v gets a score:

$$(\text{score}(v) + \epsilon) * (\text{score}(\bar{v}) + \epsilon)$$

Branch on highest scoring variable

Locally Skewed, Globally Balanced Heuristics

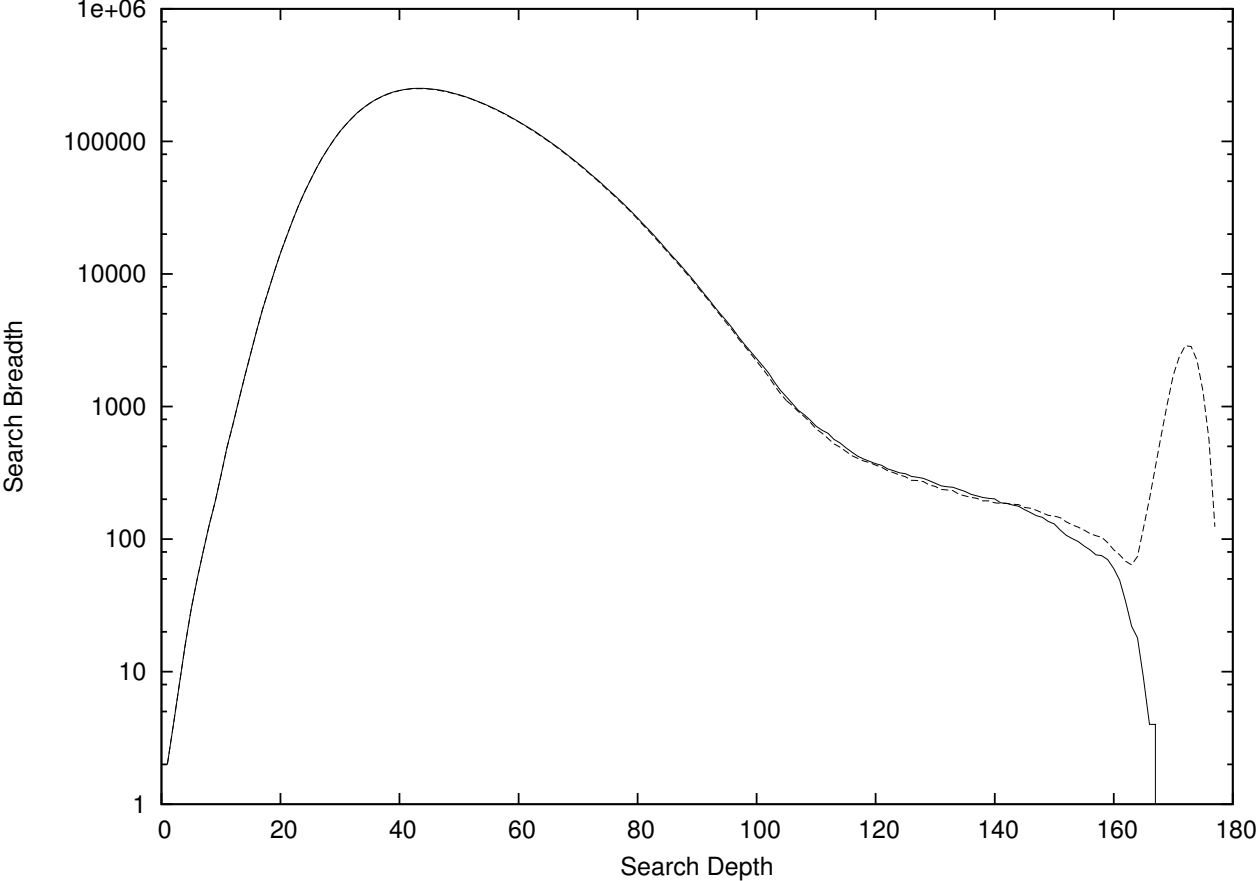


Require Pre-processing

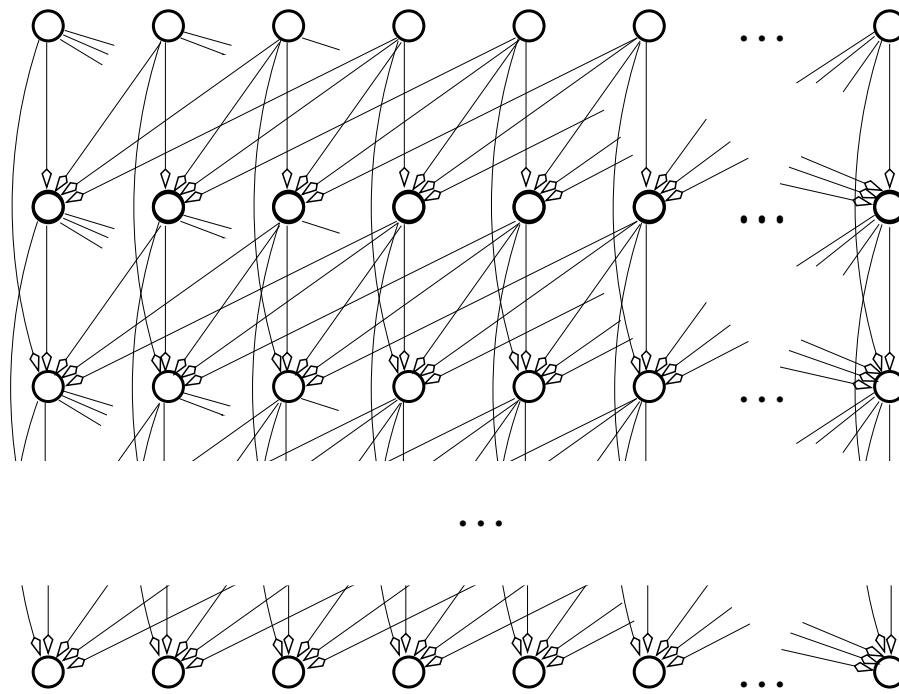
Intuitively, we desire:

- Fewer State Machines
- Smaller State Machines
- Redundancies Removed Across State Machines
- Inferences Revealed and Assigned as Early as Possible
- Safe Assignments Revealed and Assigned Early

Search Space Profile for Hard Problems



Hard Problems That Fit This Profile

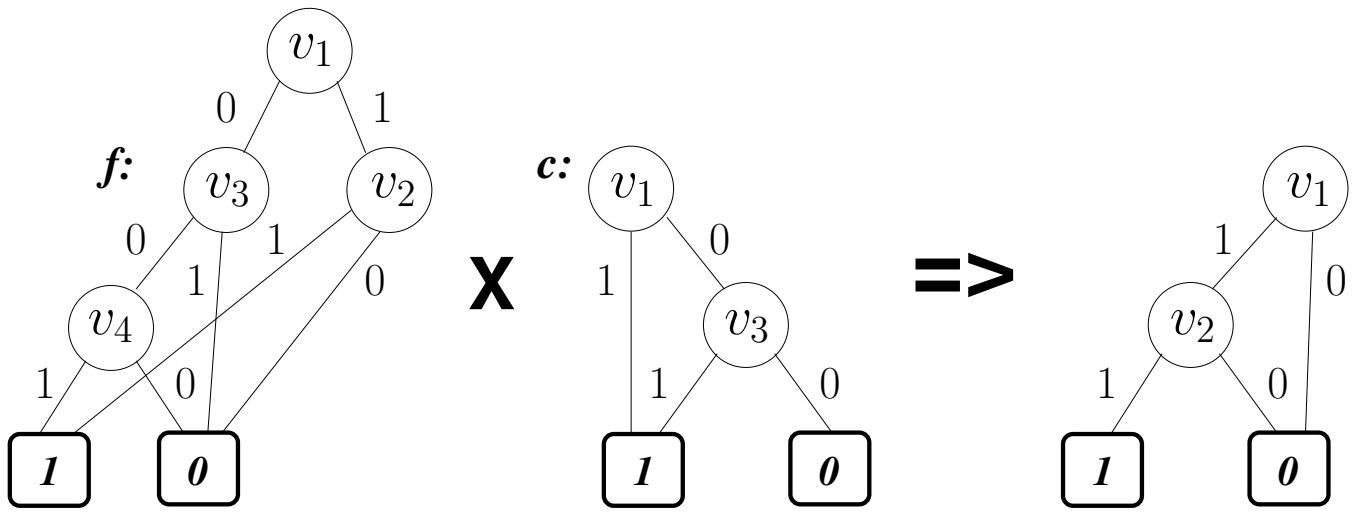


For example, Bounded Model Checking problems

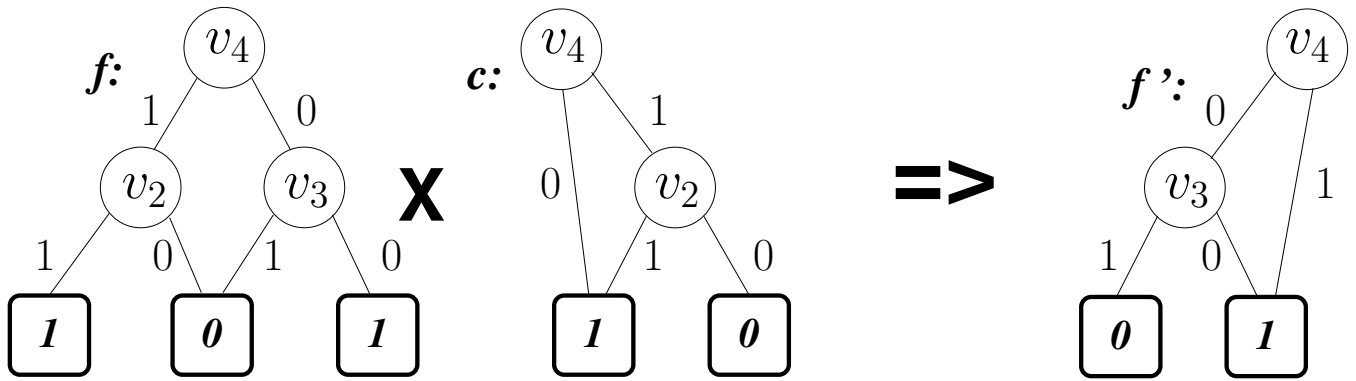
Tools for Pre-processing

- Restriction (eliminate redundancies, find inferences)
- Strengthening (find inferences missed by restriction)
- Generalized co-factor (eliminate functions)
- Cluster some functions (conjoin them)
- Existential Quantification (eliminate variables)
- Assign uninferred but safe values (reductions)
- Add uninferred and unsafe constraints (tunnel)

Restrict



Restrict

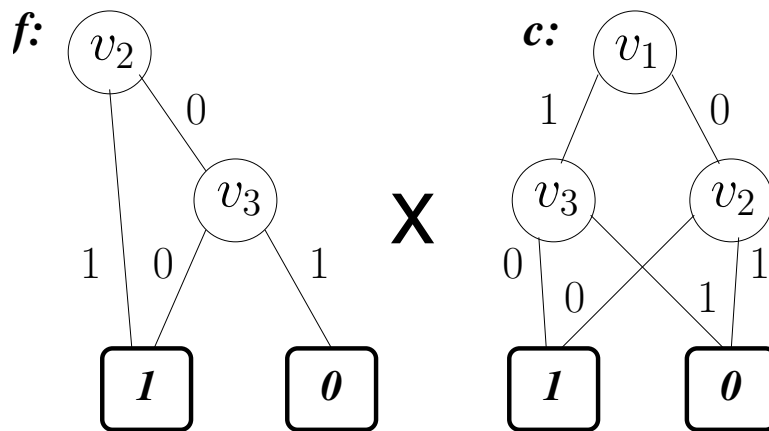


Spreading an inference from one function to another.

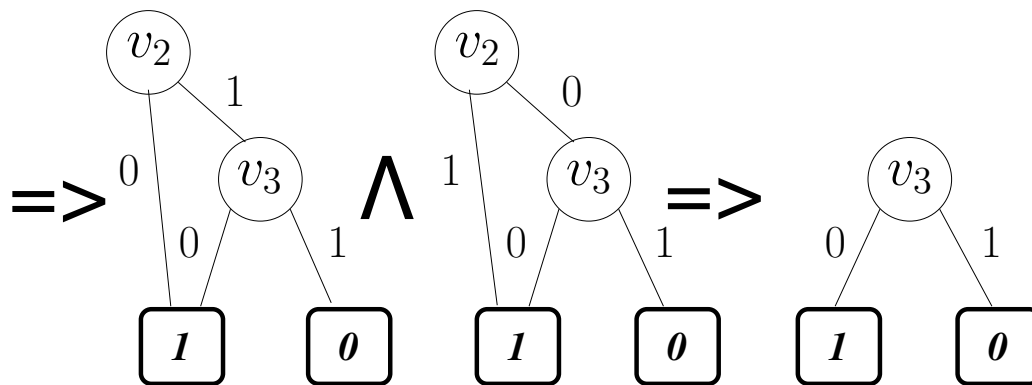
If $v_2 = 0$ in f then $v_3 = v_4 = 0$ is inferred.

Replacing f with f' , gives inference $v_4 = 0$ from c (if $v_2 = 0$)
and then inference $v_3 = 0$ from f' .

Strengthening



Existentially quantify away v_1 from f , then ...



conjoin f and c to reveal inference $v_3 = 0$.

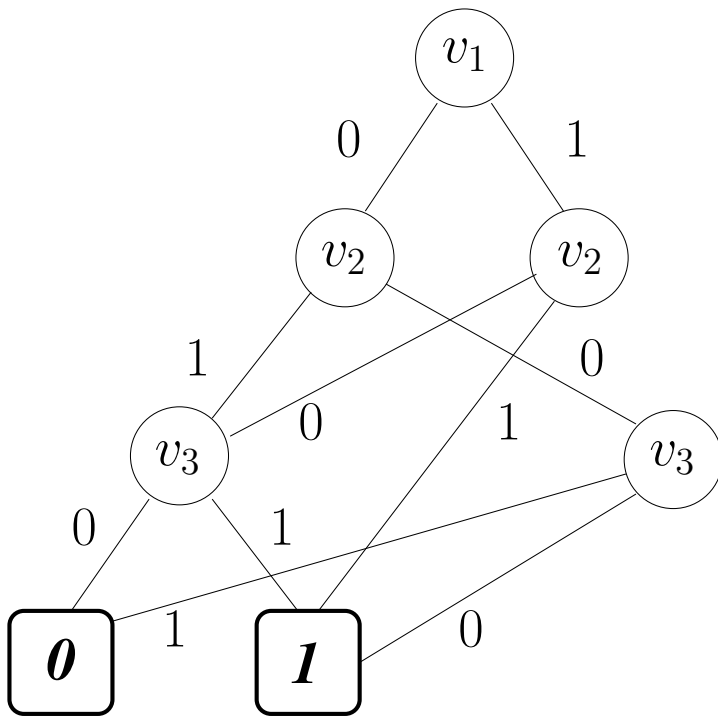
Clustering and Existential Quantification

$$\exists v(f_1 \wedge \dots \wedge f_m) \equiv (f_1 \wedge \dots \wedge f_m)|_{v=0} \vee (f_1 \wedge \dots \wedge f_m)|_{v=1}$$

But this is what we really want:

$$\exists v(f_1 \wedge \dots \wedge f_m) \equiv (f_1|_{v=0} \vee f_1|_{v=1}) \wedge \dots \wedge (f_m|_{v=0} \vee f_m|_{v=1})$$

Existential Quantification

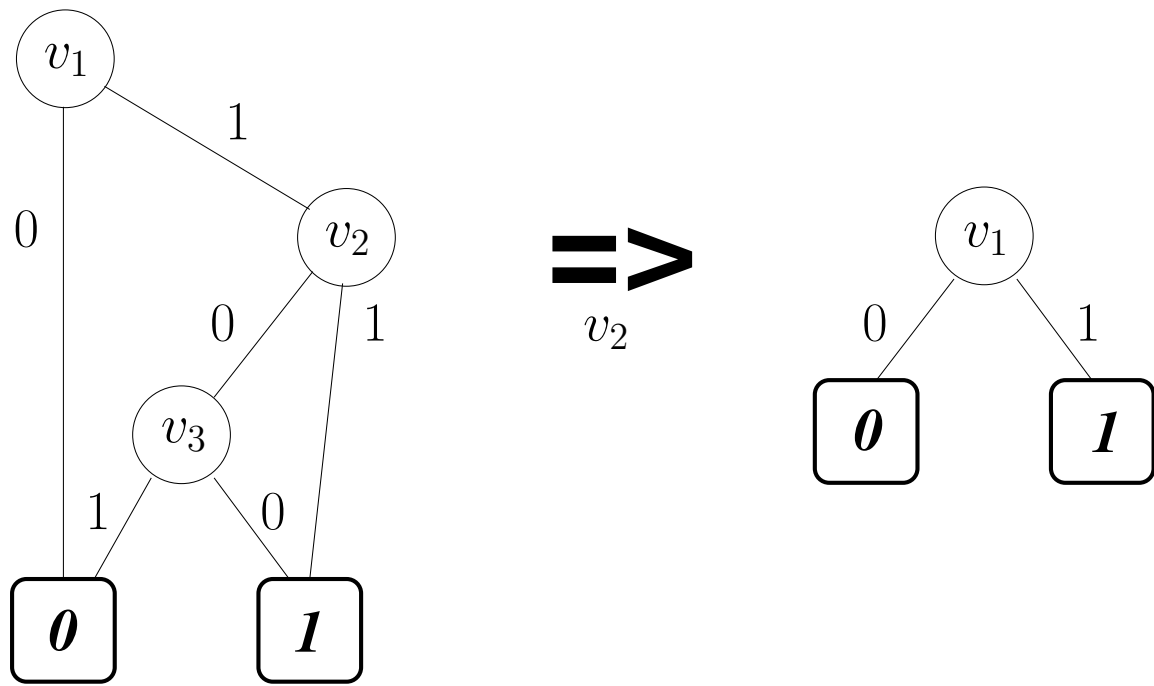


\Rightarrow
 v_3

1

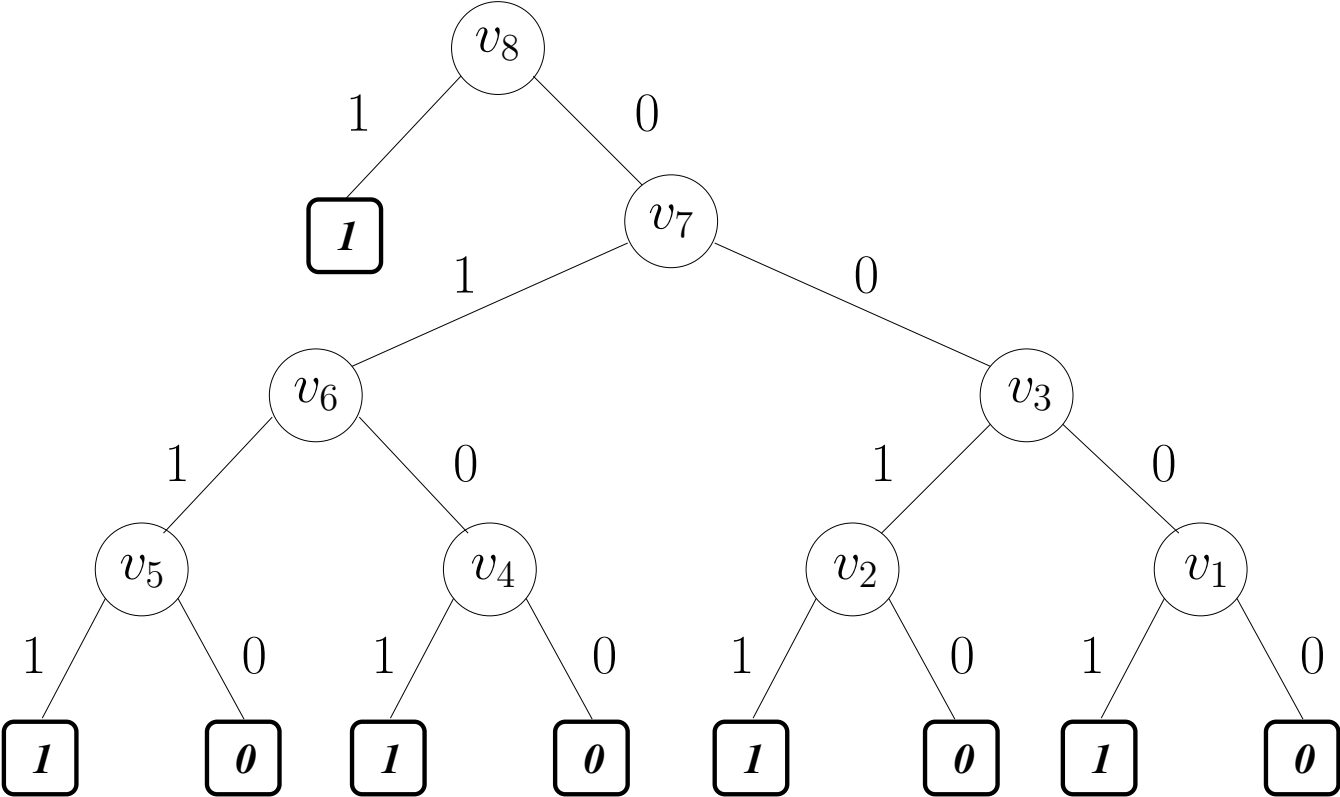
Replace f with $(f|_{v=0} \vee f|_{v=1})$

Existential Quantification



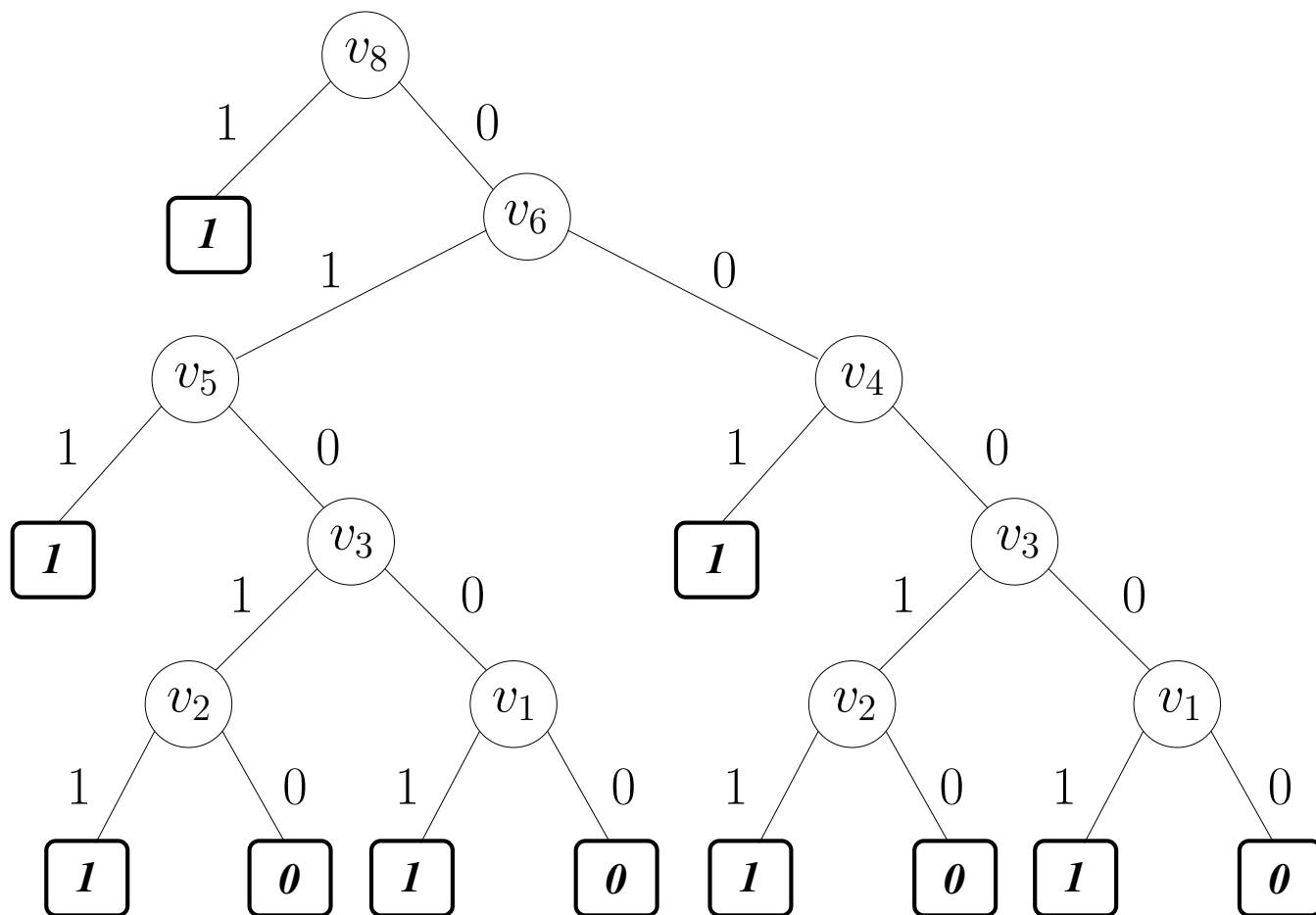
Replace f with $(f|_{v=0} \vee f|_{v=1})$

Existential Quantification



Choose v_7 to separate v_1, v_2, v_3 from v_4, v_5, v_6

Existential Quantification



But existentially quantify v_7 away and get a mess!

Splitting An Expression Is Desirable

$$\underbrace{f_1 \wedge f_2 \wedge \dots \wedge f_i}_{\text{If no variables in } f_1 \text{ to } f_i \text{ are in } f_{i+1} \text{ to } f_m} \wedge \underbrace{f_{i+1} \wedge \dots \wedge f_{m-1} \wedge f_m}_{\text{and no variables in } f_{i+1} \text{ to } f_m \text{ are in } f_1 \text{ to } f_i}$$

If no variables

in f_1 to f_i

are in f_{i+1} to f_m

and no variables

in f_{i+1} to f_m

are in f_1 to f_i

then $\phi_1 = f_1 \wedge \dots \wedge f_i$ and $\phi_2 = f_{i+1} \wedge \dots \wedge f_m$
can be solved independently

Autark Assignments Come Close

$$\underbrace{f_1 \wedge f_2 \wedge \dots \wedge f_i}_{\text{left part}} \wedge \underbrace{f_{i+1} \wedge \dots \wedge f_{m-1} \wedge f_m}_{\text{right part}}$$

If there is a subset V' of variables in f_1 to f_i and an assignment $t_{V'}$ of values to V' that satisfies f_1 to f_i

and none of the variables of V' is in f_{i+1} to f_m

then satisfy $\phi_1 = f_1 \wedge \dots \wedge f_i$ with partial assignment $t_{V'}$ and solve $\phi_2 = f_{i+1} \wedge \dots \wedge f_m$ independently

Safe Assignments

If $(f_1 \wedge \dots \wedge f_m)|_{v=0} \equiv (f_1 \wedge \dots \wedge f_m)|_{v=0} \vee (f_1 \wedge \dots \wedge f_m)|_{v=1}$
then $v = 0$ is safe

If $(f_1 \wedge \dots \wedge f_m)|_{v=1} \equiv (f_1 \wedge \dots \wedge f_m)|_{v=0} \vee (f_1 \wedge \dots \wedge f_m)|_{v=1}$
then $v = 1$ is safe

Example: $v = 1$ is a safe assignment but is not autark:

$$f = (v \vee a) \wedge (v \vee \bar{a} \vee b) \wedge (\bar{v} \vee \bar{a} \vee b \vee c) \wedge (\bar{b} \vee \bar{c}) \dots$$

Because

$$\begin{aligned} f|_{v=0} &= (a) \wedge (\bar{a} \vee b) \wedge (\bar{b} \vee \bar{c}) \rightarrow \\ f|_{v=1} &= (\bar{a} \vee b \vee c) \wedge (\bar{b} \vee \bar{c}) \end{aligned}$$

Safe Assignments: Checking isn't too bad

But conjoining functions to find safe assignments can be expensive

Luckily, the computational effort can be distributed:

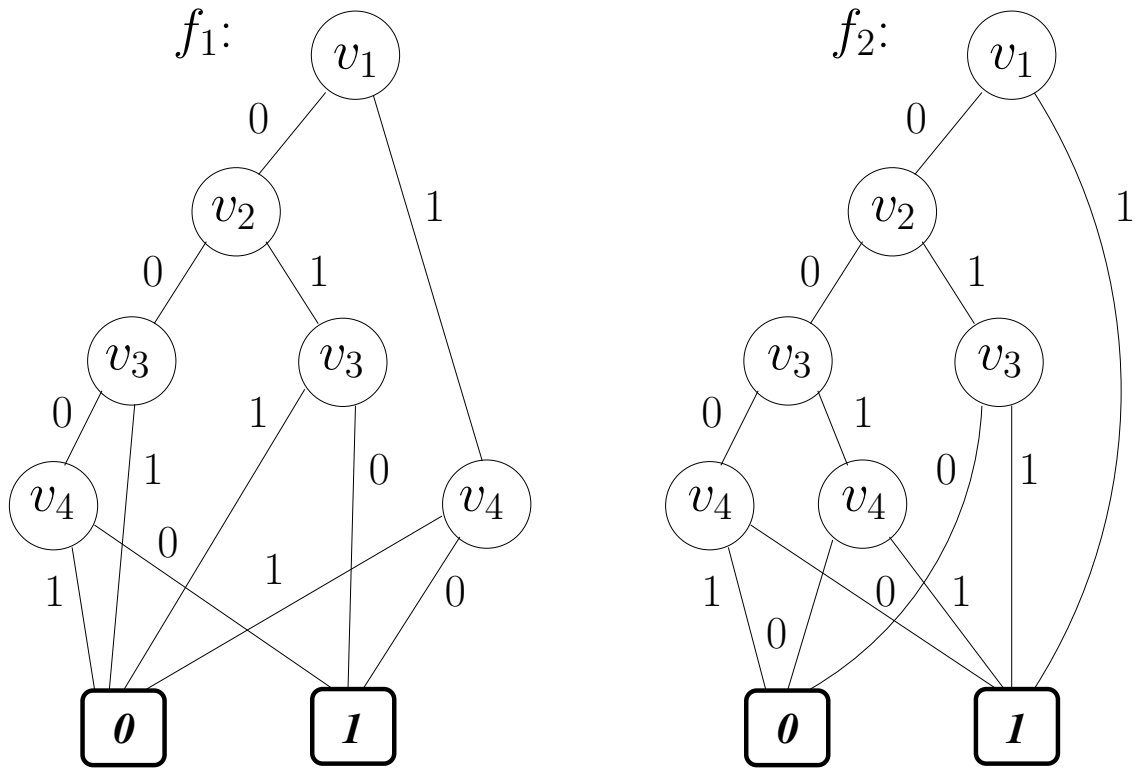
If $(\overline{f_i|_v} \wedge f_i|_{\bar{v}}) \equiv 0$ for every i such that v is in f_i then $v = 1$ is safe

If $(f_i|_v \wedge \overline{f_i|_{\bar{v}}}) \equiv 0$ for every i such that v is in f_i then $v = 0$ is safe

But the check may fail on some safe assignments

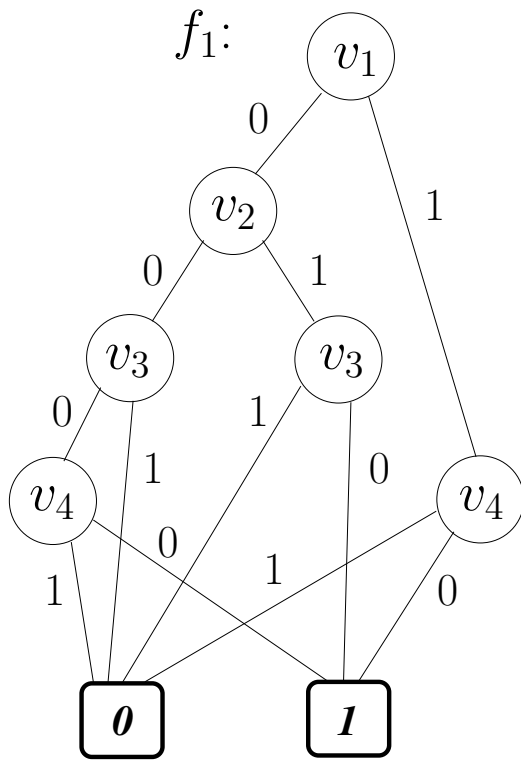
The idea extends to groups of variables

Safe Assignments: Example, Single Variable



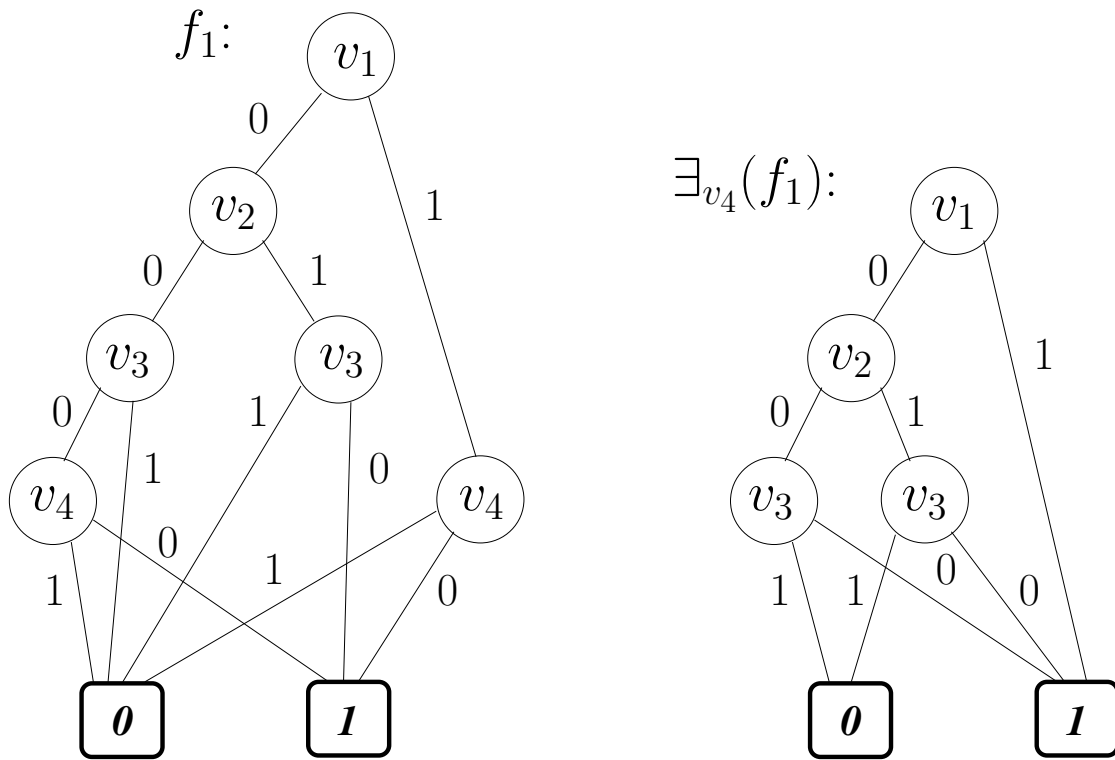
Consider two functions f_1 and f_2

Safe Assignments: Example, Single Variable



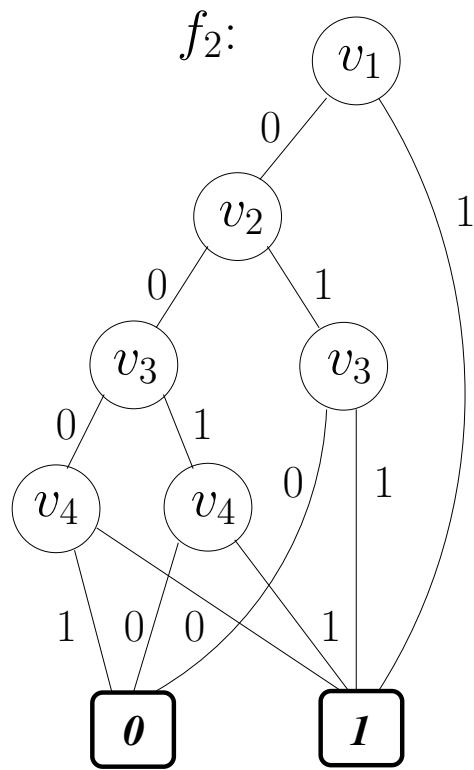
Finding safe assignment for v_4 in f_1 alone:

Safe Assignments: Example, Single Variable



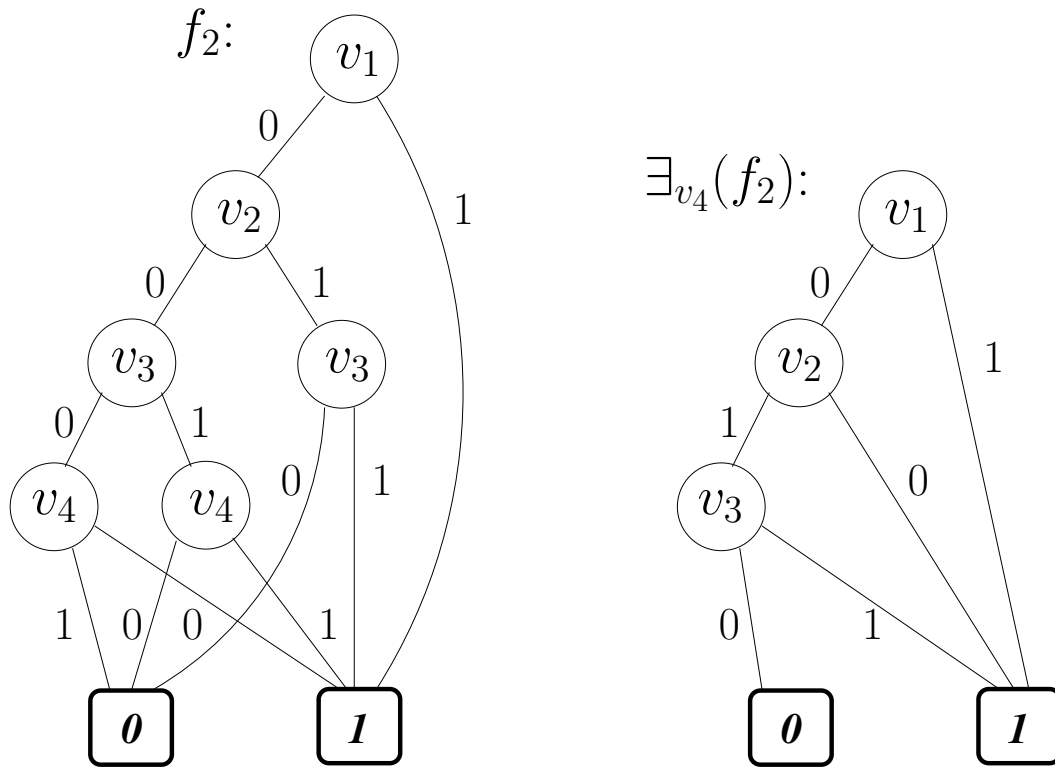
Finding safe assignment for v_4 in f_1 alone: $v_4 = 0$ only

Safe Assignments: Example, Single Variable



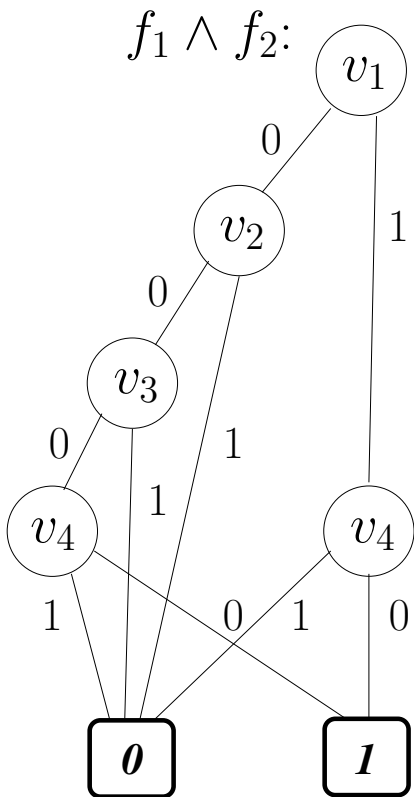
Finding safe assignment for v_4 in f_2 alone:

Safe Assignments: Example, Single Variable



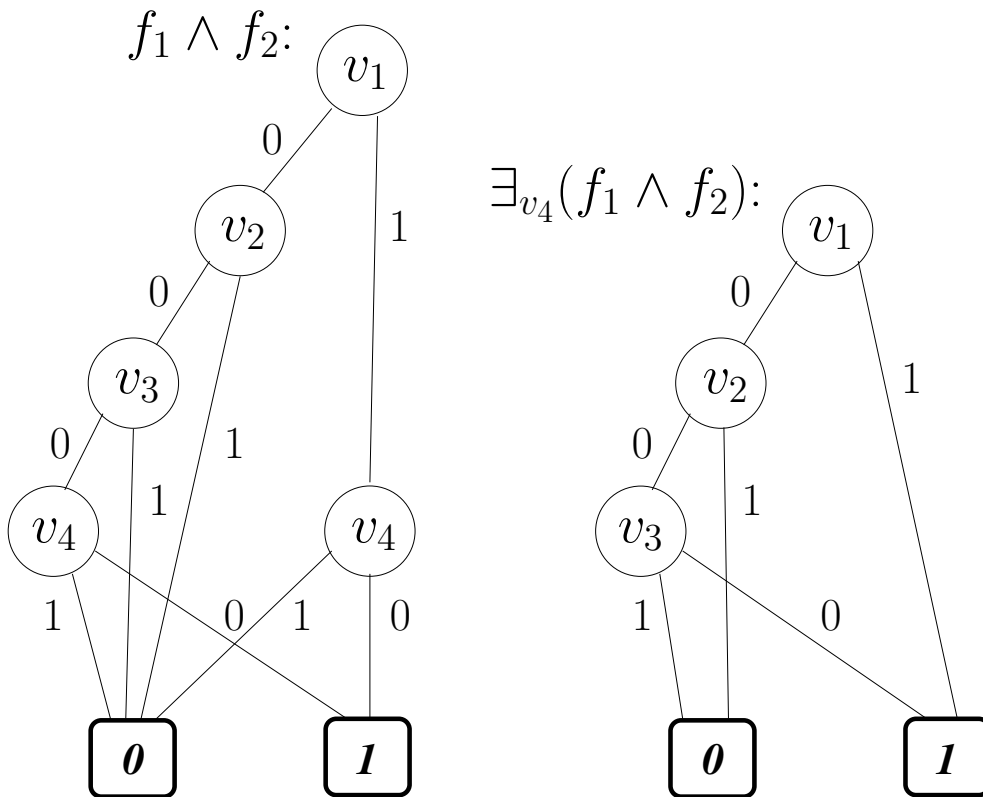
Finding safe assignment for v_4 in f_2 alone: neither one

Safe Assignments: Some Are Missed



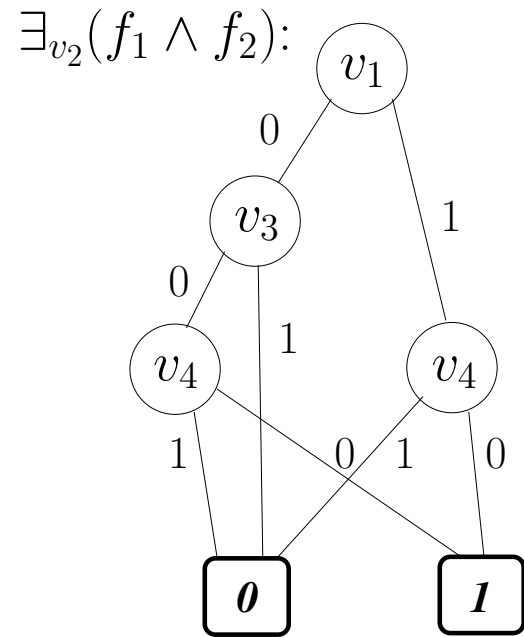
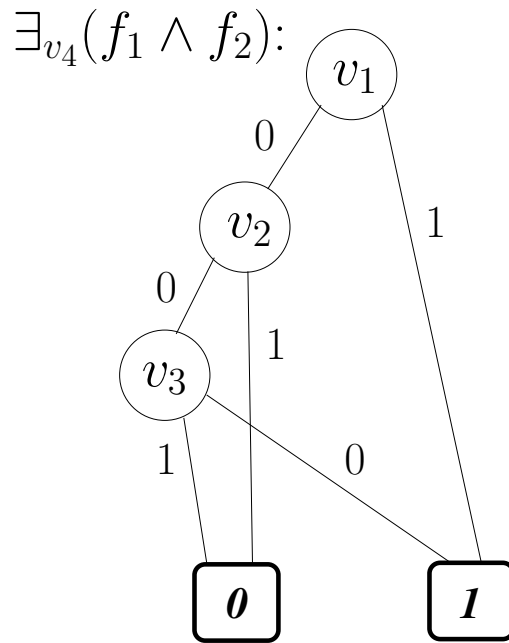
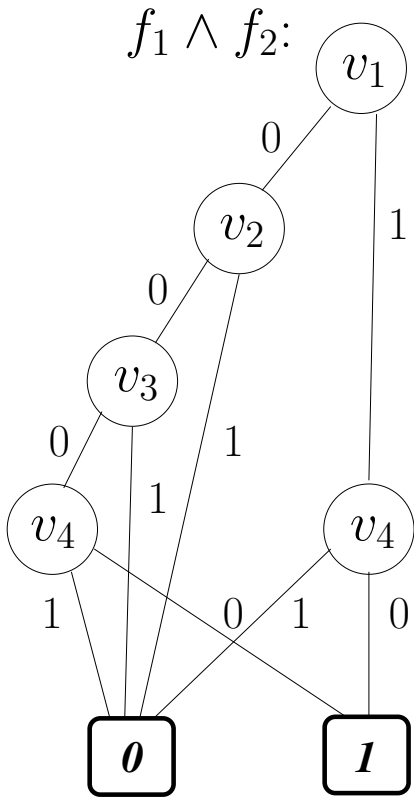
Conjoin the two functions

Safe Assignments: Some Are Missed



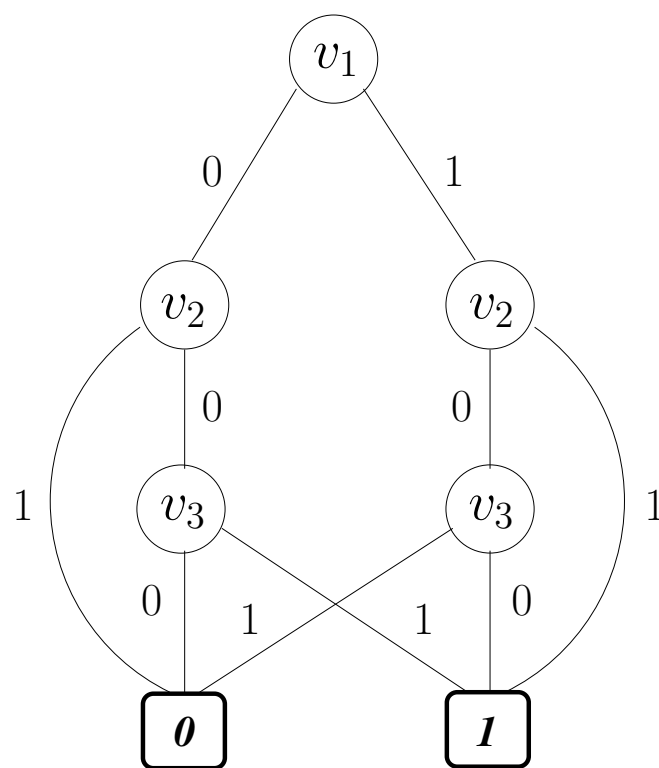
Conjoin the two functions then find $v_4 = 0$ is safe

Safe Assignments: Some Are Missed



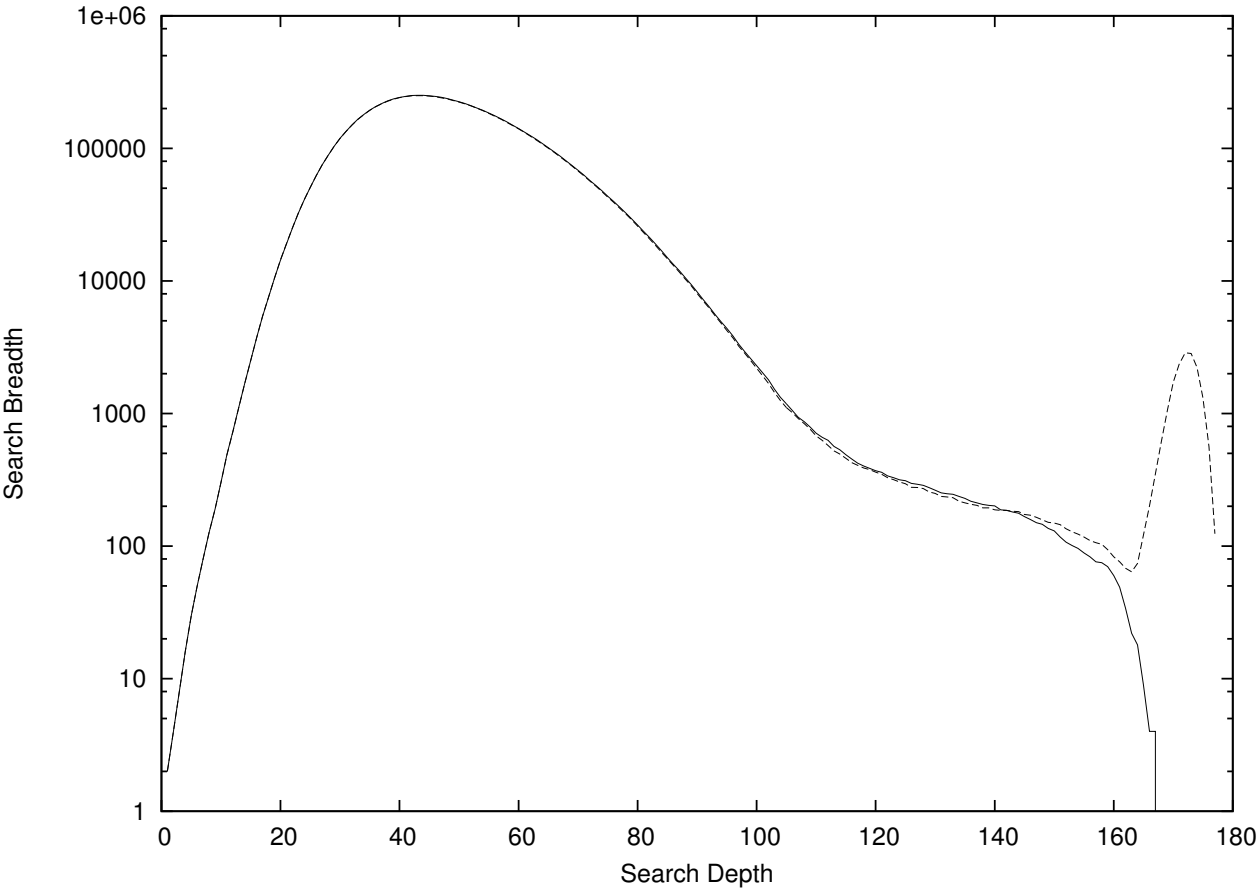
Conjoin the two functions then find $v_4 = 0$ is safe, also $v_2 = 0$ is safe

Safe Assignments: But Multiple Assignments Can Pay Off



Safe assignments $v_1 = v_2 = 1$ and $v_1 = 1, v_3 = 0$ are found as pairs only.

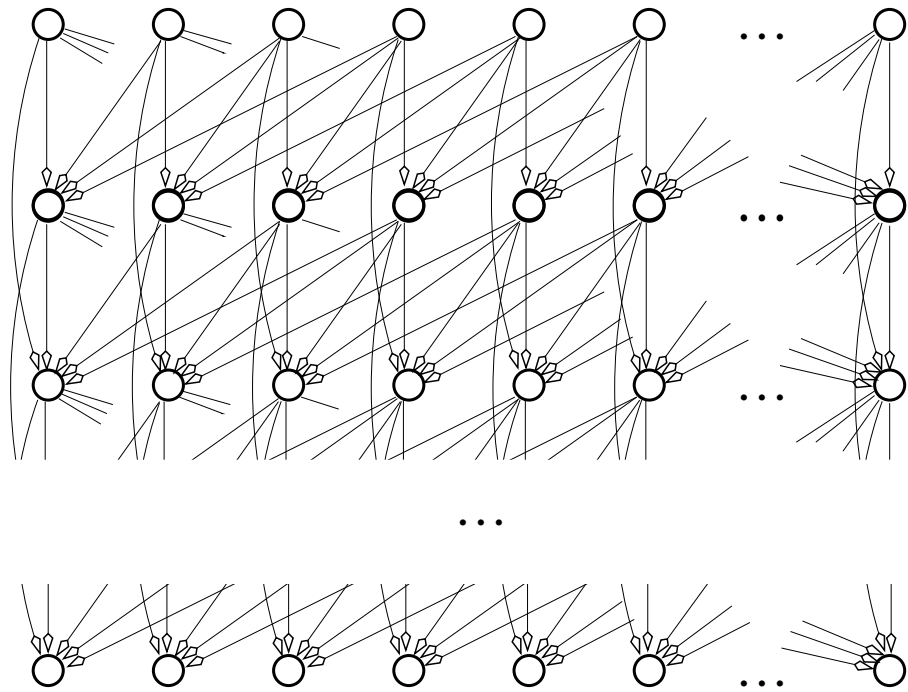
Search Space Profile for Hard Problems



Unsafe Assignments

- Guess some (uninferred) constraints based on *solution* structure in the same family
- Add those constraints initially to reduce the “hump”
- Run the search breadth-first
- When search breadth begins to decline, remove the constraints
- Solve to completion
- Possibly no solution found for a satisfiable input

Example: Van der Waerden Numbers



An ordering of the input variables is natural

Example: Van der Waerden Numbers

1 1 1 1 1 1 1 1
0 0 1 0 0 1 1 1 0 0 1 0 0
 0 1 0 0 1 0 0 1
 0 0 1 0

1010001110100100011101101000111010

2 categories, progression length 4 ($W_{max}(2, 4)$ formula)

1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
0 0 1 0 0 1 1 1 0 0 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1
 0 1 0 0 1 0 0 1 0 1 0 1 0 1 0 1 0 1
 0 0 1 0 0 1 0 1 0 1 0 1 0 1 0 1
1110111101101000101110100100001001011101000111011101110001011100001000010010111010001110111011100010111010010000111010001110111011100010111010010000100101110100010

2 categories, progression length 5 ($W_{max}(2, 5)$ formula)

Analysis of Solutions Suggests...

Conjecture:

For every $W_{max}(2, l)$ formula there exists a solution that contains at least one reflected pattern of length $W(2, l)/(2 * (l - 1))$ with the middle positioned somewhere between $W(2, l)/(l - 1)$ and $W(2, l) * (l - 2)/(l - 1)$.

From search profile:

The maximum breadth occurs near depth $W(2, l)/(2(l - 1))$.

$W(2, l) \approx l * W(2, l - 1)$, for small l anyway.

Continuing...

Add the unsafe constraints:

<u>Function Seg</u>	<u>Range</u>	<u>Meaning</u>
$(v_{-i} \vee v_{i+1}) \wedge (\bar{v}_{-i} \vee \bar{v}_{i+1})$	$0 \leq i < s/2$	force $v_{-i} \equiv \bar{v}_{i+1}$.

Retract the constraints at depth

$$(l * W(2, l - 1)) / (2(l - 1))$$

Continue without unsafe constraints until the end

Summary

State Machines

- Support function-complete look-ahead
- Efficiently support complex heuristics
- Efficiently admit special forms, e.g. cardinality constraints
- Efficiently admit backjumping, lemmas, restarts, etc.

Preprocessing

- Restrict, Existential Quantification, Strengthening help

Safe Assignments

Unsafe Assignments